Advanced Unix System Administration

Lecture 12 April 2, 2008

Steven Luo <sluo+decal@OCF.Berkeley.EDU>

- IPv6
 - IPv4 address exhaustion a looming issue (possibly as early as 2010)
 - Is already an issue in Japan and developing countries, which got less of the address allocation than existing developed countries
 - Complex CIDR network topology means that the IPv4 routing table is growing exponentially
 - IP needs to catch up with advances in routing technology and with the modern security environment

- IPv6 con't
 - Addressing: same idea as IPv4, but 128-bit addresses instead of 32-bit
 - Usual notation: 8 parts in hex separated by colons
 - By convention, high 64 bits of the IP used for the network ID, low 64 bits for the host ID
 - The (excessively) large host ID allows for stateless autoconfiguration via MAC or EUI-64 addresses
 - Router just announces prefix, clients automagically configure themselves
 - DHCPv6 available for more complex setups

- IPv6 con't
 - The original concept: hierarchical allocations and hierarchical routing
 - Regional Internet Registries get very large blocks (/16) from IANA
 - RIRs allocate /32 or larger blocks to ISPs
 - ISPs allocate /48 to organizations, /64 to known small subscribers
 - Prefixes to be allocated to facilitate prefix aggregation; would make the IPv6 routing table much smaller and easier to manage

- IPv6 con't
 - Failure of the hierarchical allocation model
 - Large organizations insisted on providerindependent (PI) address space
 - Renumbering is difficult, so PI addresses make switching ISPs easier
 - Also makes multihoming (connecting to more than one ISP) easier
 - Large endpoints can therefore now get /32 allocations on their own
 - With this allocation model, the IPv6 routing table may get to be just as large – or larger! – than the IPv4 routing table

- IPv6 con't
 - IPv4 and IPv6 can coexist, but v6 is not backwards-compatible
 - Transition mechanisms:
 - IPv6-in-IPv4 tunnels
 - Mapping the IPv4 space into the IPv6 space
 - Protocol translation mechanisms
 - Transition is going slowly, but Catch-22 involved: no one will deploy until commonly used, no one will use until commonly deployed

The Transport Layer

- User Datagram Protocol (UDP)
 - Extremely simple, but provides basically no features
 - Unreliable, message-oriented, stateless
 - Each message is packaged into a single datagram and sent over the network
 - No guarantees on delivery, order of arrival
 - UDP packet: source port (16 bits), destination port (16 bits), length (16 bits), checksum (16 bits), data

The Transport Layer

- Ports
 - In TCP and UDP, each socket is assigned a port number to identify the traffic to it
 - Kernel examines destination port to decide which process to give data to; source port used to determine destination of replies
 - By convention:
 - 1-1023 are "Well Known Ports" for services, 1024-49151 are "Registered Ports" for services,
 - 49152-65535 are dynamic ports used as source ports