Advanced Unix System Administration

Lecture 17 April 4, 2007

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- A good lot of what's on for today's networks was designed in the 1970s and 1980s for trusted networks
- This has unfortunate consequences for those of us working on a hostile Internet in the 21st century
 - Difficult to fix some of these problems without breaking backwards compatibility
 - Other problems can be fixed, but the fixes look fairly ugly

- Host-spoofing attacks
 - Various techniques, but the idea is always the same: pretend to be someone else on the network
 - If the remote service grants access based on the identity of the host, might be able to do damage
- Man-in-the-middle attacks
 - Read/modify traffic going in between hosts
 - Can be done as a router, or with a two-way host spoofing attack

- Promiscuous mode
 - Normally, an Ethernet adapter only reads traffic destined to its MAC address
 - In promiscuous mode, the adapter reads all traffic regardless of MAC address
 - On unswitched and wireless networks, this is all traffic!

- ARP cache poisoning attack
 - Recall that hosts make an ARP announcement broadcast when they plug into the network
 - By broadcasting a fake ARP announcement, we might be able to get a host to "update" its ARP cache with bad values
 - We then (hopefully) get all traffic for this IP
 - This works on switched networks too

- TCP initial sequence prediction
 - Recall the TCP three-way handshake: client SYN (with client ISN), server SYN/ACK (with server ISN, acknowledging client ISN), client ACK (acknowledging server ISN)
 - If the client can predict the server's ISN, it doesn't need to receive the server's SYN/ACK to be able to complete this connection sequence
 - This allows us to spoof being another host
 - See RFC 1948 for the classic solution

SYN flood

- To be able to finish the three-way handshake, a host (conventionally) needs to store state for each SYN it receives
- This "SYN queue" can't be allowed to grow without bound
- By filling up a host's SYN queue, we can prevent it from taking further TCP connections
 - This requires a much smaller number of packets than a straight flood
- Classic solution: TCP syncookies

- DNS cache poisoning
 - A few different ways of introducing bad entries:
 - We may be able to spoof a response from a recursive lookup
 - We could also return a fake NS record for the target domain's nameserver when the server looks up something from us
 - This bad entry then lives in the cache for the specified TTL
 - Impact similar to the ARP cache poisoning attack, except at a different layer

- Morals of the story
 - You cannot trust information you receive from the network without some verification!
 - You cannot trust the identity of the host you're talking to without some form of higher-layer authentication!
 - You don't want to allocate resources based on the initial stages of a connection
 - Segmenting your physical networks is a good idea